

A Divisibility Theorem For Factorials

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1. INTRODUCTION

Let $P(n)$ denote the largest prime factor of n , and set $\mathcal{B}(x) =$ the set of integers $n \leq x$ such that n does not divide $P(n)!$ (i.e., $P(n)$ factorial). In [1] Erdos proposed that $|\mathcal{B}(x)| = o(x)$. A solution was given in [2]. This was improved to $|\mathcal{B}(x)| = O(x \exp(-c\sqrt{\log x}))$ for some constant $c > 0$ [3]. In a sense this says something about the relative manner in which prime powers divide the integers. We propose to refine this in a different direction.

A first thought is to consider the possibility that a power of n , n^k , $k > 1$, divides $P(n)!$. But this fails since $P(n)$ divides $P(n)!$ exactly to the first power. This suggests omitting the largest prime from the factorization of n , and introducing

$$\tilde{n} = \prod_{p|n, p < P(n)} p^{v_p(n)}, \quad (1.1)$$

where $v_p(n)$ = the exponent of the highest power of p which divides n . For any function $f(x) \geq 1$ we consider

$$\mathcal{S}(f, x) = \{n \leq x, \tilde{n}^{f(x)} \text{ does not divide } P(n)!\}, \quad (1.2)$$

and wish to characterize those $f(x)$ such that

$$|\mathcal{S}(f, x)| = o(x). \quad (1.3)$$

($|\mathcal{S}|$ denotes the number of elements in the set \mathcal{S} .)

DEFINITION 1.1. The function $f(x)$ belongs to the lower class \mathcal{L} if (1.3) holds. Otherwise $f(x) \in \mathcal{U}$, the upper class.

Remark. $f(x) \in \mathcal{U}$ is equivalent to the existence of a constant $c > 0$ and a sequence of $x_i \rightarrow \infty$ such that $|\mathcal{S}(f, x_i)| \geq c x_i$ for all i .

We pass to a logarithmic scale and introduce

$$\varepsilon(f, x) = \varepsilon(x) = \frac{\log f(x)}{\log x}. \quad (1.4)$$

Our main objective is to establish

THEOREM 1.1. A necessary and sufficient condition for $f(x) \in \mathcal{L}$ is that

$$\lim_{x \rightarrow \infty} \varepsilon(f, x) = 0. \quad (1.5)$$

(Hence when (1.5) fails $f(x) \in \mathcal{U}$.)

2. THE NECESSITY OF (1.5)

Assuming that (1.5) fails, we wish to show that $f(x) \in \mathcal{U}$. Then

$$\overline{\lim}_{x \rightarrow \infty} \varepsilon(f, x) > 0, \quad (2.1)$$

so that on some sequence $x_i \rightarrow \infty$, $\varepsilon(f, x_i) > c > 0$. Thus it suffices to show that there exists a constant $d > 0$ such that

$$|\mathcal{S}(f, x_i)| > d x_i. \quad (2.2)$$

An integer $n \in \mathcal{S}(f, x)$ if

$$f(x) \nu_p(n) > \nu_p(P(n)!) \quad (2.3)$$

for some prime $p \mid n$, $p < P(n)$. Since $\nu_p(n) \geq 1$ and $\nu_p(P(n)!) \leq P(n)$, (2.3) is implied by

$$f(x) \geq P(n). \quad (2.4)$$

The set of $n \leq x$ satisfying (2.4) is contained in $\mathcal{S}(f, x)$ and it suffices to show that for $x = x_i$, (2.4) is satisfied by more than $d x_i$ values of $n \leq x_i$.

This is an immediate consequence of the following

LEMMA 2.1. *For any fixed constant $c > 0$, there exists a constant $d > 0$, depending on c , such that the number of $n \leq x$ for which*

$$P(n) \leq x^c \quad (2.5)$$

is greater than $d x$, for all sufficiently large x .

(The following proof was suggested by P. Erdos. In a later section we'll consider the actual density, $\delta(c)$, of these integers, which provides another proof as well as a method for calculating $\delta(c)$.)

Proof. It clearly suffices to consider the cases where $c = \frac{1}{k}$, k an integer ≥ 2 . Choose ε , $0 < \varepsilon < \frac{1}{k+1}$, and consider the integers $n \leq x$ which are divisible by an integer of the form

$$m = p_1 p_2 \cdots p_k, \quad x^{(1-\varepsilon)\frac{1}{k}} < p_i \leq x^{\frac{1}{k}}, \quad i = 1, \dots, k, \quad (2.6)$$

where the p_i are primes. Such an integer is of the form tm where $1 \leq t < x^\varepsilon$, so that all the prime factors of t are less than $x^\varepsilon < x^{1/k}$. Further, our choice of ε implies that $\varepsilon < (1 - \varepsilon)\frac{1}{k}$, so that the representations in the form tm are unique. Thus, the integers divisible by those of the form (2.6) satisfy the condition (2.5) and have a count of at least

$$\begin{aligned} \frac{1}{k!} \frac{x}{2} \sum_{x^{\frac{1-\varepsilon}{k}} < p_i \leq x^{\frac{1}{k}}} \frac{1}{p_1 \cdots p_k} &= \frac{x}{2k!} \left(\sum_{x^{\frac{1-\varepsilon}{k}} < p \leq x^{\frac{1}{k}}} \frac{1}{p} \right)^k \\ &= \frac{x}{2k!} \left(\log \frac{1}{1 - \varepsilon} + o(1) \right)^k \end{aligned}$$

and the lemma follows.

3. THE SUFFICIENCY OF (1.5)

Our next objective is to show that (1.5) implies that the number of $n \leq x$ such that (2.3) holds, for some prime $p < P(n)$ which divides n , is $o(x)$. Here we must replace $\mathcal{S}(f, x)$ by an appropriate larger set. Also, without loss of generality we may assume that $f(x) \rightarrow \infty$ as $x \rightarrow \infty$. Since $\nu_p(P(n)!) \geq [\frac{P(n)}{p}] > \frac{P(n)}{2p}$, (2.3) implies

$$\nu_p(n) f(x) > \frac{P(n)}{2p}, \quad (3.1)$$

for some prime p which divides n .

The set of $n \leq x$ such that $v_p(n) \geq T$, for some prime p , is at most

$$\begin{aligned} x \sum_p \left\{ \frac{1}{p^T} + \frac{1}{p^{T+1}} + \cdots \right\} &= x \sum_p \frac{1}{p^T} \left(1 - \frac{1}{p} \right)^{-1} \\ &\leq 2x \sum_{n=2}^{\infty} \frac{1}{n^T} = O \left(x \int_2^{\infty} \frac{du}{u^T} \right) = O \left(\frac{x}{T} \right), \end{aligned}$$

which is $o(x)$ for $T \rightarrow \infty$ however slowly. Thus the set of interest is contained in the $n \leq x$ such that

$$Tf(x) > \frac{P(n)}{2p}, \quad (3.2)$$

for some prime $p < P(n)$ which divides n . But such an n is divisible by a prime p such that

$$P(n) > p > \frac{P(n)}{2Tf(x)}. \quad (3.3)$$

We propose to show that the number of such n is $o(x)$.

Since $P(n)$ is the largest prime dividing n , these n are relatively prime to all primes less than x and greater than $P(n)$. We then have two cases:

Case I. $P(n) < x^{\varepsilon_1(x)}$ where $\varepsilon_1(x) \rightarrow 0$ is chosen later.

Then the number of $n \leq x$ which are prime to all primes $p > P(n)$ is at most the number of $n \leq x$ which are prime to all $p > x^{\varepsilon_1(x)}$. Next choose an $\varepsilon_2(x) > \varepsilon_1(x)$, which also tends to 0, but more slowly than ε_1 , that is, so that $\varepsilon_1(x)/\varepsilon_2(x) \rightarrow 0$ as $x \rightarrow \infty$. Hence the integers under consideration are prime to all p in the interval $x^{\varepsilon_1(x)} < p < x^{\varepsilon_2(x)}$. It is a well known sieving result [4] that this number is

$$O \left(x \prod_{x^{\varepsilon_1} < p \leq x^{\varepsilon_2}} \left(1 - \frac{1}{p} \right) \right) = O \left(x \frac{\varepsilon_1(x)}{\varepsilon_2(x)} \right),$$

which is $o(x)$.

Case II. $P(n) > x^{\varepsilon_1(x)}$, where $\varepsilon_1(x) \rightarrow 0$ (as adjusted again later).

We choose $\varepsilon_1(x) \rightarrow 0$ so that $\varepsilon(x)/\varepsilon_1(x) \rightarrow 0$ and $x^{\varepsilon_1}/f(x) = x^{\varepsilon_1 - \varepsilon} \rightarrow \infty$. Note that this implies that $x^{\varepsilon_1}/Tf(x) \rightarrow \infty$ if $T \rightarrow \infty$ sufficiently slowly.

Then the $n \leq x$ under consideration are divisible by a prime $q = P(n)$, and by another prime p , where

$$\frac{q}{2Tf(x)} < p < q. \quad (3.4)$$

The number of such n is at most

$$x \sum_{x^{\varepsilon_1} < q \leq x} \frac{1}{q} \sum_{\substack{q \\ \frac{q}{2Tf(x)} < p < q}} \frac{1}{p}. \quad (3.5)$$

Using the fact that

$$\sum_{p \leq z} \frac{1}{p} = \log \log z + C + O\left(\frac{1}{\log z}\right)$$

(for some constant C), we have

$$\sum_{\substack{q \\ \frac{q}{2Tf(x)} < p < q}} \frac{1}{p} = \log \log q - \log \log \frac{q}{2Tf(x)} + O\left(\left(\log \frac{q}{2Tf(x)}\right)^{-1}\right). \quad (3.6)$$

The first two terms on the right of (3.6) equal

$$\log \frac{\log q}{\log(q/2Tf(x))} = O\left(\frac{\log 2Tf(x)}{\log q}\right),$$

and the O term

$$\frac{1}{\log(q/2Tf(x))} = O\left(\frac{\log 2Tf(x)}{\log q}\right).$$

Hence (3.5) equals

$$O\left(x \sum_{x^{\varepsilon_1} < q \leq x} \frac{\log 2Tf(x)}{q \log q}\right) = O\left(x \frac{\log 2Tf(x)}{\varepsilon_1 \log x}\right).$$

Also choosing T so that $\log T = O(\varepsilon \log x)$, the above equals $O(x(\varepsilon/\varepsilon_1)) = o(x)$, since $\varepsilon/\varepsilon_1 \rightarrow 0$ as $x \rightarrow \infty$. (Note that this uses the fact that $f(x) \rightarrow \infty$.)

This completes the proof of Theorem 1.1.

4. ON THE VALUE OF $\delta(c)$

We return now to a detailed look at the value of the density $\delta(c)$ (whose existence has not as yet been demonstrated here). Recall that this is the density of integers such that (2.5) holds. That is, those integers $n \leq x$, all of whose prime factors are less than or equal to x^c .

Clearly we may assume that $c < 1$. Also, such n are precisely those which are relatively prime to all primes p such that

$$x^c < p \leq x. \quad (4.1)$$

Choose the integer $k \geq 1$ where

$$\frac{1}{k+1} \leq c < \frac{1}{k}. \quad (4.2)$$

The density of the complement of the set of integers described above (if it exists),

$$\bar{\delta}(c) = 1 - \delta(c) \quad (4.3)$$

is given by

$$\bar{\delta}_k(c) = \bar{\delta}(c) = \sum_{i=1}^k (-1)^{i+1} \lim_{x \rightarrow \infty} \sum_{\substack{x^c < p_{j_1} < \dots < p_{j_i} \leq x \\ \prod_{\lambda=1}^i p_{j_\lambda} \leq x}} (p_{j_1} \cdots p_{j_i})^{-1}. \quad (4.4)$$

Setting $G_i(c)$ equal to the last summation in (4.4) and $F_i(c)$ equal to the limit of this sum, we have

$$\bar{\delta}(c) = \sum_{i=1}^k (-1)^{i+1} F_i(c). \quad (4.5)$$

LEMMA 4.1. *For any fixed c , $0 < c < 1$ (recall (4.2)), for $1 \leq m \leq k$,*

$$F_m(c) = \int \cdots \int_{\mathcal{D}_m} \prod_{i=1}^m \frac{dx_i}{x_i}, \quad (4.6)$$

where the domain $\mathcal{D}_m = \mathcal{D}_m(c)$ is given by

$$\mathcal{D}_m = \left\{ c < x_1 < x_2 < \cdots < x_m, \sum_{i=1}^m x_i \leq 1 \right\}. \quad (4.7)$$

(Note that for $m > k$, $mc > 1$, and $F_m(c) = 0$.)

Proof. By the Prime Number Theorem [4], if p_i is the i th prime $p_i = (i \log i)(1 + o(1))$, then

$$G_m(c) = \sum_{\mathcal{A}_m} \left[\prod_{j=1}^m (i_j \log i_j)^{-1} (1 + o(1)) \right], \quad (4.8)$$

where the domain \mathcal{R}_m is given by

$$\mathcal{R}_m = \left(c < \left(\frac{\log i_1}{\log x} \right) (1 + o(1)) < \cdots < \left(\frac{\log i_m}{\log m} \right) (1 + o(1)) \leq 1 \right. \\ \left. \sum_{j=1}^m \left(\frac{\log i_j}{\log x} \right) (1 + o(1)) \leq 1 \right). \quad (4.9)$$

Writing

$$\prod_{j=1}^m (i_j \log i_j)^{-1} = \prod_{j=1}^m \left(\frac{\log i_j}{\log x} \right)^{-1} (i_j \log x)^{-1}$$

we see that the right side of (4.9) is a Riemann sum for the integral in (4.6). (Note that the $(1 + o(1))$ factors are easily disposed of.)

LEMMA 4.2. *For c and k as specified in Lemma 4.1, and $1 \leq m \leq k$,*

$$F_m(c) = \int_c^{\frac{1}{m}} F_{m-1} \left(\frac{u}{1-u} \right) \frac{du}{u}. \quad (4.10)$$

(Note that $F_0(c) \equiv 1$ is consistent with the above.)

Proof. Clearly (4.10) holds for $m = 1$, and we assume $m \geq 2$. From (4.6)

$$F_m(c) = \int_c^{\frac{1}{m}} \frac{dx_1}{x_1} \int_{\mathcal{K}_{m-1}} \cdots \int \prod_{i=2}^m \frac{dx_i}{x_i}, \quad (4.11)$$

where \mathcal{K}_{m-1} is the domain

$$\mathcal{K}_{m-1} = \mathcal{K}_{m-1}(x_1) = \left[x_1 < x_2 < \cdots < x_m, \sum_{i=2}^m x_i \leq 1 - x_1 \right]. \quad (4.12)$$

Introducing new variables $y_i = x_i/(1 - x_1)$, $i = 2, \dots, m$, the domain of the inner integer becomes

$$\frac{x_1}{1 - x_1} < y_2 < \cdots < y_m, \sum_{i=2}^m y_i < 1,$$

and

$$\prod_{i=2}^m \frac{dx_i}{x_i} = \prod_{i=2}^m \frac{dy_i}{y_i}$$

so that (4.11) yields (4.10).

LEMMA 4.3. For any $c > 0$, $\frac{1}{k+1} \leq c < \frac{1}{k}$, and all $m = 1, \dots, k$, $F_m(c)$ is differentiable at c , and

$$c \frac{d}{dc} F_m(c) = -F_{m-1} \left(\frac{c}{1-c} \right). \quad (4.13)$$

Proof. Since $F_0(c) \equiv 1$ we have $F_1(c) = -\log c$, which has derivative $-\frac{1}{c}$ in the interval $\frac{1}{2} \leq c < 1$. Proceeding by induction we assume the lemma for $m-1$ where $2 \leq m \leq k$. Since the mapping $u \rightarrow \frac{u}{1-u}$ takes the interval $[\frac{1}{k+1}, \frac{1}{k})$ onto $[\frac{1}{k}, \frac{1}{k-1})$, F_{m-1} is certainly continuous at $\frac{c}{1-c}$. Hence $F_m(c)$ has a derivative at c and (4.13) follows from (4.10).

LEMMA 4.4. For $0 < c < 1$, $\frac{1}{k+1} \leq c < \frac{1}{k}$, we have

$$-c \frac{d}{dc} \bar{\delta}_k(c) = 1 - \bar{\delta}_{k-1} \left(\frac{c}{1-c} \right). \quad (4.14)$$

Proof. Applying (4.5) and (4.13) for $m = 1, \dots, k$, we obtain

$$\begin{aligned} -c \frac{d}{dc} \bar{\delta}_k(c) &= \sum_{m=1}^k (-1)^m c \frac{d}{dc} F_m(c) \\ &= 1 - \sum_{m=2}^k (-1)^m F_{m-1} \left(\frac{c}{1-c} \right), \end{aligned}$$

and (4.14) follows.

COROLLARY. The density $\delta(c)$ satisfies

$$c \frac{d}{dc} \delta_k(c) = \delta_{k-1} \left(\frac{c}{1-c} \right). \quad (4.15)$$

Another Proof of Lemma 2.1. Consider first the case of $\frac{1}{2} < c \leq 1$. Then

$$\bar{\delta}(c) = \bar{\delta}_1(c) = F_1(c) = \log \frac{1}{c} < \log 2 < 1.$$

Assuming that we have a real c such that $\bar{\delta}_k(c) = 1$, it follows that $0 < c \leq \frac{1}{2}$. Setting

$$c^* = \limsup_{\delta(c)=1} c, \quad (4.16)$$

for all $c < c^*$ we have $\bar{\delta}(c) = 1$. Then if $\frac{1}{k+1} < c^* \leq \frac{1}{k}$, we have $k \geq 2$, and there is an interval of c values $\frac{1}{k+1} < c < c^*$ such that $\bar{\delta}(c) = 1$ for all of these. Then $\frac{d}{dc}\bar{\delta}(c) = 0$, and (4.14) yields that $\bar{\delta}_{k-1}(\frac{c}{1-c}) = 1$. But $c^* \leq \frac{1}{k} < \frac{c}{1-c}$, which contradicts (4.16).

Turning next to the calculation of $\delta(c)$, integrating (4.15) yields

$$\delta_k\left(\frac{1}{k}\right) - \delta_k(c) = \int_c^{\frac{1}{k}} \frac{1}{u} \delta_{k-1}\left(\frac{u}{1-u}\right) du. \quad (4.17)$$

Making a change of variable this becomes

$$\delta_k(c) = \delta_k\left(\frac{1}{k}\right) - \int_{\frac{c}{1-c}}^{\frac{1}{k-1}} \frac{1}{u(u+1)} \delta_{k-1}(u) du. \quad (4.18)$$

This last provides a natural numerical algorithm for calculating $\delta(c)$.

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